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## Chapter Boolean Algebra

## Lecture =07 Karnaugh Map (خريطة كارنوف)

## 7-1. The Karnaugh Map

- The Karnaugh map (K-map) is similar to a truth table.
$\square$ The K-map is an array of cells (حصفوفة من الخالاي) in which each cell represent a binary value of the input variables.
The K-map is a tool for simplifying combinational logic with two, three, four, and five variables.
For 3 variables, 8 cells are required $\left(2^{3}\right)$, and for 4 variables: $2^{4}=16$ cells.

The map shown is for three variables labeled $A, B$, and $C$.

- Each cell represents one possible product term (مصطأح جدائي).
- Each cell differs from an adjacent (المجاورة) cell by only one variable.



## Karnaugh map

- Cells are usually labeled using 0 's and 1's to represent the variable and its complement.
- The numbers are entered in Gray code, to force adjacent cells to be different by only one variable.
- Ones are read as the true variable and zeros are read as the complemented variable.



## Example 7-1

Read the terms for the yellow cells.


## The 3-variable Karnaugh map

$\square$ The 3-variable K-maps is any array of eight cells, as shown in Fig. (a).
$\square$ Fig. (b) shows the standard product terms that are represented by each cell in the 3 -variable K-map.


| 0 | $\bar{A} \bar{B} \bar{C}$ | $\bar{A} \bar{B} C$ |
| :---: | :---: | :---: |
| 01 | $\bar{A} B \bar{C}$ | $\bar{A} B C$ |
| 11 | $A B \bar{C}$ | $A B C$ |
| 10 | $A \bar{B} \bar{C}$ | $A \bar{B} C$ |

(a)
(b)

## The 4-variable Karnaugh map

$\square$ The 4-variable K-maps is any array of sixteen cells, as shown in Fig. (a).
$\square$ Fig. (b) shows the standard product terms that are represented by each cell in the 4 -variable K-map.


## 7-2. Karnaugh Map SOP Minimization

## 7-2-1. Mapping a standard SOP Expression

For an SOP expression in standard form, a 1 is placed on the K-map for each product term in the expression.

- For example, for the product term $A \bar{B} C$, a $\mathbf{1}$ goes in the 101 cell on a 3variables map.
- The cells that do not have a 1 are the cells for which the expression is $\mathbf{0}$.
- The mapping process:

Step 1. Determine the binary value of each product sum in the standard SOP expression.
Step 2. As each product term is evaluated, place a 1 on a K-map in the cell having the same value as the product term.


Map the following standard SOP expression on a K-map:

$$
\bar{A} \bar{B} C+\bar{A} B \bar{C}+A B \bar{C}+A B C
$$

Evaluate the expression as show below.

$$
\begin{aligned}
& \bar{A} \bar{B} C+\bar{A} B \bar{C}+A B \bar{C}+A B C \\
& 001 \quad 010 \quad 110
\end{aligned}
$$

Place a 1 on the 3-variable K-map in Fig. for each standard product term in the expression.


Map the following standard SOP expression on a K-map:

$$
\bar{A} \bar{B} C D+\bar{A} B \bar{C} \bar{D}+A B \bar{C} D+A B C D+A B \bar{C} \bar{D}+\overline{A B} \bar{C} D+A \bar{B} C \bar{D}
$$

Evaluate the expression as show below.

$$
\left.\begin{array}{l}
\bar{A} \bar{B} C D+\bar{A} B \bar{C} \bar{D}+A B \bar{C} D+A B C D+A B \bar{C} \bar{D}+\bar{A} \bar{B} \bar{C} D+A \bar{B} C \bar{D} \\
0011 \\
0100 \\
1101 \\
1111 \\
\hline 1100
\end{array}\right) 001 \quad 1010
$$

Place a 1 on the 4 -variable K-map in Fig. for each standard product term in the expression.


## 7-2-2. Karnaugh Map Simplification of SOP Expression

- After an SOP expression has been mapped, a minimum SOP expression is obtained by grouping the 1 s and determining the minimum SOP expression from map.


## Grouping the 1 s .

- You can group 1s on the K-map by enclosing (إحاطة) those adjacent (المتجاورة) cells containing 1 s .
- Each group is either an square or rectangle shape.


##  <br> Group 1s in each of the K-maps:


(a)

(b)

(d)

## Solution

Wrap-around adjacency
(c)



(a)

(b)

(c)

(d)
d)
$\square$ K-maps can simplify combinational logic by grouping cells and eliminating (إقصاء) variables that change.

EAIIID7-5 Group the 1 's on the map and read the minimum logic.


## Solution

1. Group the 1's into two overlapping groups as indicated.
2. Read each group by eliminating (إزالة/إقصاء) any variable that changes across a boundary (الحد).
3. The vertical group is read $\bar{A} \bar{C}$.
4. The horizontal group is read $A B$.

$$
X=\bar{A} \bar{C}+\bar{A} B
$$

$\square$ A 4-variable map has an adjacent cell on each of its four boundaries. Example 7-6

Group the 1 's on the map and read the minimum logic. Solution

1. Group the 1 's into two separate groups as indicated.
2. Read each group by eliminating any variable that changes across a boundary.
3. The upper (yellow) group is read as $\bar{A} \bar{D}$.
4. The lower (green) group is read as $A D$.

$$
X=\bar{A} \bar{D}+A D
$$



## 7-2-3. Mapping Directly from a Truth Table

- An example of a Boolean expression and its truth table representation is in Fig.
- Notice in truth table that the output $X$ is $\mathbf{1}$ for four different input variable combinations.
- The 1s in the output column of the truth table are mapped directly onto a Kmap into the cells corresponding to the values of the associated input variable combinations, as shown in Fig.



## 7-2-4. Don't Care Conditions

$\square$ Sometimes a situation arises in which some input variable combinations are not allowed.
[ These combinations can be treated (يتّ التعامل) as "don't care" (غير مهم/ مهمل) terms with respect to (بالنسبة) لـ

- That is, for these "don't care" terms either a $\mathbf{1}$ or a $\mathbf{0}$ may be assigned to the output.
$\square$ For each "don't care" term, an X is placed in the cell.
$\square$ When grouping the 1s, the Xs can be treated as 1 s to make a larger grouping or as 0s if they cannot be used to advantage (للاستفادة/زيادة المنفعة).
- The larger a group, the simpler the resulting term will be.
$\square$ The truth table in Fig. (a) describes a logic function that has a 1 output only when the BCD code for 7,8 , or 9 is present on the inputs.
- If the "don't cares" are used as 1 s , the resulting expression for the function is $A+B C D$, as indicated in Fig.(b).
- If the "don't cares" are not used as 1 s , the resulting expression is $A \bar{B} \bar{C}+\bar{A} B C D$
- So, this we can see the advantage of using "don't care" terms to get the simplest expression.

| Inputs |  |  |  | Output |
| :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{A}$ | $\boldsymbol{B}$ | $\boldsymbol{C}$ | $\boldsymbol{D}$ | $\boldsymbol{Y}$ |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | X |
| 1 | 0 | 1 | 1 | X |
| 1 | 1 | 0 | 0 | X |
| 1 | 1 | 0 | 1 | X |
| 1 | 1 | 1 | 0 | X |
| 1 | 1 | 1 | 1 | X |


(b) Without "don't cares" $Y=A \bar{B} \bar{C}+\bar{A} B C D$ With "don't cares" $Y=A+B C D$

## Selected Key Terms

| Variable | A symbol used to represent a logical quantity that can have a value of <br> 1 or 0, usually designated by an italic letter. |
| :--- | :--- |
| Complement | The inverse or opposite of a number. In Boolean algebra, the inverse <br> function, expressed with a bar over the variable. |
| Sum term | The Boolean sum of two or more literals equivalent to an OR <br> operation. |
| Product term | The Boolean product of two or more literals equivalent to an AND <br> operation. |
| Sum-of-products <br> (SOP) | A form of Boolean expression that is basically the ORing of ANDed <br> terms. |
| Product of sums <br> (POS) | An arrangement of cells representing combinations of literals in a <br> Boolean expression and used for systematic simplification of the <br> expression. |
| "Don't care" | A combination of input literals that cannot occur and can be used as a <br> 1 or a on a Karnaugh map for simplification. |

## True/False Quiz.

1. Variable, complement, and literal are all terms used in Boolean algebra.
2. Addition in Boolean algebra is equivalent to the NOR function.
3. Multiplication in Boolean algebra is equivalent to the AND function.
4. The commutative law, associative law, and distributive law are all laws in Boolean algebra.
5. The complement of 0 is 0 itself.
6. When a Boolean variable is multiplied by its complement, the result is the variable.
7. "The complement of a product of variables is equal to the sum of the complements of each variable" is a statement of DeMorgan's theorem.
8. SOP means sum-of-products.
9. Karnaugh maps can be used to simplify Boolean expressions.
10. A 3 -variable Karnaugh map has six cells.
11. T 2. F
12. T
13. T
14. F
15. F
16. T
17. T
18. T
19. F

## SELF-TEST-1

1. The associative law for addition is normally written as
a. $A+B=B+A$
b. $(A+B)+C=A+(B+C)$
c. $A B=B A$
d. $A+A B=A$
2. The Boolean equation $A B+A C=A(B+C)$ illustrates
a. the distribution law
b. the commutative law
c. the associative law
d. DeMorgan's theorem
3. The Boolean expression $A \cdot 1$ is equal to
a. $\boldsymbol{A}$
b. $B$
c. 0
d. 1
4. The Boolean expression $A+1$ is equal to
a. $A$
b. $B$
c. 0
5. The Boolean equation $A B+A C=A(B+C)$ illustrates
a. the distribution law
b. the commutative law
c. the associative law
d. DeMorgan's theorem
6. A Boolean expression that is in standard SOP form is
a. the minimum logic expression
b. contains only one product term
c. has every variable in the domain in every term
d. none of the above
7. Adjacent cells on a Karnaugh map differ from each other by
a. one variable
b. two variables
c. three variables
d. answer depends on the size of the map
8. The minimum expression that can be read from the Karnaugh map shown is
a. $X=A$
b. $X=A$
c. $X=B$
d. $X=B$

9. The minimum expression that can be read from the Karnaugh map shown is
a. $X=A$
b. $X=A$
c. $X=B$
d. $X=\boldsymbol{B}$

## SELE-TEST-2

1. A variable is a symbol in Boolean algebra used to represent
(a) data
(b) a condition
(c) an action
(d) answers (a), (b), and (c)
2. The Boolean expression $A+B+C$ is
(a) a sum term
(b) a literal term
(c) an inverse term
(d) a product term
3. The Boolean expression $\overline{A B C D}$ is
(a) a sum term
(b) a literal term
(c) an inverse term
(d) a product term
4. The domain of the expression $A \bar{B} C D+A \bar{B}+\bar{C} D+B$ is
(a) $A$ and $D$
(b) B only
(c) $A, B, C$, and $D$
(d) none of these
5. According to the associative law of addition,
(a) $A+B=B+A$
(b) $A=A+A$
(c) $(A+B)+C=A+(B+C)$
(d) $A+0=A$
6. According to commutative law of multiplication,
(a) $A B=B A$
(b) $A=A A$
(c) $(A B) C=A(B C)$
(d) $A 0=A$
7. According to the distributive law,
(a) $A(B+C)=A B+A C$
(b) $A(B C)=A B C$
(c) $A(A+1)=A$
(d) $A+A B=A$
8. (d)
9. (a)
10. (d)
11. (c)
12. (c)
13. (a)
14. (a)
15. Which one of the following is not a valid rule of Boolean algebra?
(a) $A+1=1$
(b) $A=\bar{A}$
(c) $A A=A$
(d) $A+0=A$
16. Which of the following rules states that if one input of an AND gate is always 1 , the output is equal to the other input?
(a) $A+1=1$
(b) $A+A=A$
(c) $A \cdot A=A$
(d) $A \cdot 1=A$
17. According to DeMorgan's theorems, the complement of a product of variables is equal to
(a) the complement of the sum
(b) the sum of the complements
(c) the product of the complements
(d) answers (a), (b), and (c)
18. The Boolean expression $X=(A+B)(C+D)$ represents
(a) two ORs ANDed together
(b) two ANDs ORed together
(c) A 4-input AND gate
(d) a 4-input OR gate
19. An example of a sum-of-products expression is
(a) $A+B(C+D)$
(b) $\bar{A} B+A \bar{C}+A \bar{B} C$
(c) $(\bar{A}+B+C)(A+\bar{B}+C)$
(d) both answers (a) and (b)
20. An example of a product-of-sums expression is
(a) $A(B+C)+A \bar{C}$
(b) $(A+B)(\bar{A}+B+\bar{C})$
(c) $\bar{A}+\bar{B}+B C$
(d) both answers (a) and (b)
21. An example of a standard SOP expression is
(a) $\bar{A} B+A \bar{B} C+A B \bar{D}$
(b) $A \bar{B} C+A \bar{C} D$
(c) $A \bar{B}+\bar{A} B+A B$
(d) $A \bar{B} C \bar{D}+\bar{A} B+\bar{A}$
22. (b) $\quad$ 9. (d) $\quad 10$. (b) 11 . (a) $\quad 12$. (b) $\quad 13$. (b) $\quad 14$. (c)

## Problems \& Solutions

Find the values of the variables that make each product term 1 and each sum term 0 .
(a) $A B$
(b) $A \bar{B} C$
(c) $A+B$
(d) $\bar{A}+B+\bar{C}$
(e) $\bar{A}+\bar{B}+C$
(f) $\bar{A}+\hat{B}$
(g) $A \bar{B} C$
(a) $A B=1$ when $A=B=1$
(b) $A \bar{B} C=1$ when $A=1, B=0, C=1$
(c) $A+B=0$ when $A=0, B=0$
(d) $\bar{A}+B+\bar{C}=0$ when $A=1, B=0, C=1$
(e) $\bar{A}+\bar{B}+C=0$ when $A=1, B=1, C=0$
(f) $\bar{A}+B=0$ when $A=1, B=0$
(g) $A \bar{B} \bar{C}=1$ when $A=1, B=0, C=0$

Find the value of $\boldsymbol{X}$ for all possible values of the variables.
(a) $X=(A+B) C+B$
(b) $X=(\overline{A+B}) C$
(c) $X=A \bar{B} C+A B$
(d) $X=(A+B)(\bar{A}+B)$
(e) $X=(A+B C)(\bar{B}+\bar{C})$

## Sol.

(a) $\quad X=(A+B) C+B$
(b) $\quad X=(\overline{A+B}) C$
(c) $X=A \bar{B} C+A B$


| $A$ | $B$ | $C$ | $A \bar{B} C$ | $A B$ | $X$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 | 1 |

$$
\text { (d) } \quad X=(A+B)(\bar{A}+B)
$$

| $A$ | $B$ | $A+B$ | $\bar{A}+B$ | $X$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 |

(e) $\quad X=(A+B C)(\bar{B}+\bar{C})$

| $A$ | $B$ | $C$ | $A+B C$ | $\bar{B}+\bar{C}$ | $X$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 |

(a) $\overline{A+\bar{B}}$
(b) $\overline{\bar{A} B}$
(c) $\overline{A+B+C}$
(d) $\overline{A B C}$
(e) $\overline{A(B+C)}$
(f) $\overline{A B}+\overline{C D}$
(g) $\overline{A B+C D}$
(h) $\overline{(A+\bar{B})(\bar{C}+D)}$
(a) $\overline{A+\bar{B}}=\bar{A} \bar{B}=\bar{A} \boldsymbol{B}$
(b) $\bar{A} B=A+\bar{B}=\boldsymbol{A}+\overline{\boldsymbol{B}}$
(c) $\overline{A+B+C}=\overline{\boldsymbol{A}} \overline{\boldsymbol{B}} \overline{\boldsymbol{C}}$
(d) $\overline{A B C}=\overline{\boldsymbol{A}}+\overline{\boldsymbol{B}}+\overline{\boldsymbol{C}}$
(e) $\overline{A(B+C)}=\bar{A}+\overline{(B+C)}=\overline{\boldsymbol{A}}+\overline{\boldsymbol{B}} \overline{\boldsymbol{C}}$
(f) $\overline{A B}+\overline{C D}=\bar{A}+\bar{B}+\bar{C}+\bar{D}$
(g) $\overline{A B+C D}=\overline{(A B)(C D)}=(\overline{\boldsymbol{A}}+\overline{\boldsymbol{B}})(\overline{\boldsymbol{C}}+\overline{\boldsymbol{D}})$
(h)

$$
\overline{(A+\bar{B})(\bar{C}+D)}=\overline{A+\bar{B}}+\overline{\bar{C}+D}=\overline{\boldsymbol{A}} \boldsymbol{B}+\boldsymbol{C} \overline{\boldsymbol{D}}
$$

Prob. $4-4$
Write the Boolean expression for each of the logic gates in Fig.

(a)
$A->0-x$
(b)

(c)
${ }_{C}^{B} \longrightarrow-X$
(d)
(a) $\quad A B=X$
(c) $A+B=X$
(b) $\bar{A}=X$
(d) $A+B+C=X$
$* * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * * *$

## Prob. 4.5

Write the Boolean expression for each of the logic circuits in Fig.

(a)

(b)


(d)
(a) $\stackrel{x}{x}=A B C D$

(c) $\mathrm{X}=\overline{\overline{A B}}$

(b) $X=A B+C$

(d) $X=(A+B) C$


PO1.4-6 Draw the logic circuit represented by each of the following expressions:
$\begin{array}{llll}\text { (a) } A+B+C & \text { (b) } A B C & \text { (c) } A B+C & \text { (d) } A B+C D\end{array}$

(a)

(c)

(b)

(d)

Construct a truth table for each of the following Boolean
expressions:
(a) $A+B$
(b) $A B$
(c) $A B+B C$
(d) $(A+B) C$
(e) $(A+B)(\bar{B}+C)$
(a) $A+B$

| $A$ | $B$ | $X$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

(c) $\quad X=A B+B C$

| $A$ | $B$ | $C$ | $X$ |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |


| $A$ | $B$ | $C$ | $X$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

(e) $\quad X=(A+B)(\bar{B}+C)$

| $A$ | $B$ | $C$ | $A+B$ | $\bar{B}+C$ | $X$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 | 1 |

Using Boolean algebra, simplify the following expressions:
(a) $B D+B(D+E)+\bar{D}(D+F)$
(b) $\bar{A} \bar{B} C+\overline{(A+B+\bar{C})}+\bar{A} \bar{B} \bar{C} D$
(c) $(B+B C)(B+\bar{B} C)(B+D)$
(d) $A B C D+A B(\overline{C D})+(\overline{A B}) C D$
(e) $A B C[A B+\bar{C}(B C+A C)]$

$$
\begin{equation*}
B D+B(D+E)+\bar{D}(D+F)=B D+B D+B E+\bar{D} D+\overline{D F} \tag{a}
\end{equation*}
$$

$$
=B D+B E+0+\bar{D} F=\boldsymbol{B} \boldsymbol{D}+\boldsymbol{B} \boldsymbol{E}+\overline{\boldsymbol{D}} \boldsymbol{F}
$$

(b) $\bar{A} \bar{B} C+(\overline{A+B+\bar{C}})+\overline{A B} \bar{C} D=\overline{A B} C+\overline{A B} C+\overline{A B} \bar{C} D=\bar{A} \bar{B} C+\overline{A B} \bar{C} D$

$$
=\bar{A} \bar{B}(C+\bar{C} D)=\bar{A} \bar{B}(C+D)=\overline{A B} C+\overline{A B} D
$$

(c) $\quad(B+B C)(B+\bar{B} C)(B+D)=B(1+C)(B+C)(B+D)$
$=B(B+C)(B+D)=(B B+B C)(B+D)=(B+B C)(B+D)$
$=B(1+C)(B+D)=B(B+D)=B B+B D=B+B D=B(1+D)=B$

$$
\begin{equation*}
A B C D+A B(\overline{C D})+(\overline{A B}) C D=A B C D+A B(\bar{C}+\bar{D})+(\bar{A}+\bar{B}) C D \tag{d}
\end{equation*}
$$

$=A B C D+A B \bar{C}+A B \bar{D}+\bar{A} C D+\bar{B} C D$
$=C D(A B+\bar{A}+\bar{B})+A B \bar{C}+A B \bar{D}=C D(B+\bar{A}+\bar{B})+A B \bar{C}+A B \bar{D}$
$=C D(1+\bar{A})+A B \bar{C}+A B \bar{D}=C D+A B \bar{C}+A B \bar{D}=C D+A B(\overline{C D})=\boldsymbol{C D}+A \boldsymbol{B}$
$A B C[A B+\bar{C}(B C+A C)]=A B A B C+A B C \bar{C}(B C+A C)$
$A B C+0(B C+A C)=A B C$

Convert the following expressions to sum-of-product (SOP) forms:
(a) $A B+C D(A \bar{B}+C D)$
(b) $A B(\bar{B} \bar{C}+B D)$
(c) $A+B[A C+(B+\bar{C}) D]$
(a) $A B+C D(A B+C D)=A B+A B C D+C D C D=A B+A B C D+C D$ $=A B(A \bar{B}+1) C D=A B+C D$
(b) $A B(\overline{B C}+B D)=A B \overline{B C}+A B B D=0+A B D=A B D$
(c) $A+B[A C+(B+\bar{C}) D]=A+A B C+(B+\bar{C}) B D$

$$
=A+A B C+B D+B \bar{C} D=A(1+B C)+B D+B \bar{C} D=A+B D(1+\bar{C})
$$

$$
=A+B D
$$

Develop a truth table for each of the SOP expressions:
$\begin{array}{ll}\text { (a) } \bar{A} B+A B \bar{C}+\bar{A} \bar{C}+A \bar{B} C & \text { (b) } \bar{X}+Y \bar{Z}+W Z+X \bar{Y} Z\end{array}$
(a) $\bar{A} B+A B \bar{C}+\bar{A} \bar{C}+\overline{A B C}=\bar{A} B C+\bar{A} B \bar{C}+A B \bar{C}+\overline{A B C}+\overline{A B} C$
(b)

$$
\begin{aligned}
\bar{X}+Y \bar{Z} & +W Z+X \bar{Y} Z=\bar{W} \bar{X} \bar{Y} \bar{Z}+\bar{W} \bar{X} \bar{Y} Z+\bar{W} X Y \bar{Z}+\bar{W} \bar{X} Y Z \\
& +\bar{W} X \bar{Y} Z+\bar{W} X Y \bar{Z}+W X \bar{Y} \bar{Z}+W \bar{X} \bar{Y} Z \\
& +W \bar{X} Y \bar{Z}+W \bar{X} Y Z+W X Y \bar{Y}+W X Y \bar{Z}+W X Y Z
\end{aligned}
$$

| $A$ | $B$ | $C$ | $X$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 0 |


| $W$ | $X$ | $Y$ | $Z$ | $Q$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |

Develop a truth table for each of the standard POS expressions:
(a) $(A+B)(A+C)(A+B+C)$
(b) $(A+\bar{B})(A+\bar{B}+\bar{C})(B+C+\bar{D})(\bar{A}+B+\bar{C}+D)$
(a)

| $A$ | $B$ | $C$ | $X$ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

(b)

| $A$ | $B$ | $C$ | $D$ | $X$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |

P101.412 For each truth table, derive a standard SOP and a standard POS expression.

(a)

$$
\begin{aligned}
& X=\bar{A} \bar{B} C+A \bar{B} \bar{C}+A \bar{B} C+A B C \\
& X=(A+B+C)(A+\bar{B}+C)(A+\bar{B}+\bar{C})(\bar{A}+\bar{B}+C)
\end{aligned}
$$

(b) $\quad X=A B \bar{C}+A \bar{B} C+A B C$

$$
X=(A+B+C)(A+B+\bar{C})(A+\bar{B}+C)(A+\bar{B}+\bar{C})(\bar{A}+B+C)
$$

$$
\begin{aligned}
X= & \overline{A B C} \bar{C}+\overline{A B C} D+\overline{A B} C D+\bar{A} B \bar{C} D+\bar{A} B C \bar{D}+A \overline{B C} D+A B \bar{C} \bar{D} \\
X= & (A+B+\bar{C}+D)(A+\bar{B}+C+D)(A+\bar{B}+\bar{C}+\bar{D})(\bar{A}+B+C+D)(\bar{A}+B+\bar{C}+D) \\
& (\bar{A}+B+\bar{C}+\bar{D})(\bar{A}+\bar{B}+C+\bar{D})(\bar{A}+\bar{B}+\bar{C}+D)(\bar{A}+\bar{B}+\bar{C}+\bar{D})
\end{aligned}
$$

Draw a 3-variable Karnaugh map and label each cell according to its binary value.

## Sol.

| 00 | 000 | 001 |
| :---: | :---: | :---: |
| 01 | 010 | 011 |
| 11 | 110 | 111 |
| 10 | 100 | 101 |

****************************************************************
Proln $4-14$ Write the standard product term for each cell in a 3-variable Karnaugh map.

| 00 | $\bar{A} \bar{B} \bar{C}$ | $\bar{A} \overline{B C}$ |
| :---: | :---: | :---: |
| 01 | $\bar{A} B \bar{C}$ | $\bar{A} B C$ |
| 11 | $A B C$ | $A B C$ |
| 10 | $A \bar{B} \bar{C}$ | $A \overline{B C}$ |

Given the Karnaugh maps in Figure. Write the resulting minimum SOP expression for each.

(a)

(b)

(c)

(d)

The resulting minimum product term for each group is shown in Figure. The minimum SOP expressions for each of the Karnaugh maps in the figure are
(a) $A B+B C+\bar{A} \bar{B} \bar{C}$
(b) $\bar{B}+\bar{A} \bar{C}+A C$
(c) $\bar{A} B+\bar{A} \bar{C}+A \bar{B} D$
(d) $\bar{D}+A \bar{B} C+B \bar{C}$

Expand each expression to a standard SOP form:
(a) $A B+A \bar{B} C+A B C$
(b) $A+B C$
(c) $A \bar{B} \bar{C} D+A C \bar{D}+B \bar{C} D+\bar{A} B C \bar{D}$
(d) $A \bar{B}+A \bar{B} \bar{C} D+C D+B \bar{C} D+A B C D$
(a) $A B+A \bar{B} C+A B C=A B(C+\bar{C})+A \bar{B} C+A B C$

$$
\begin{aligned}
& =A B C+A B \bar{B}+A \bar{B} C+A B C \\
& =A B C+A \bar{B} C+A B \bar{C}
\end{aligned}
$$

(b) $A+B C=A(B+\bar{B})(C+\bar{C})+(\bar{A}+A) B C=(A B+A \bar{B})(C+\bar{C})+(\bar{A}+A) B C$

$$
=A B C+A B \bar{C}+A \bar{B} C+A \bar{B} \bar{C}+\bar{A} B C+A B C
$$

$$
=A B C+A \bar{B} \bar{C}+A \bar{B} C+A \bar{B} \bar{C}+\bar{A} B C
$$

(c) $A \bar{B} \bar{C} D+A C \bar{D}+B \bar{C} D+\overline{A B C} \bar{D}$

$$
\begin{aligned}
& =A \bar{B} \bar{C} D+A(B+\bar{B}) C D+(A+\bar{A}) B \bar{C} D+\bar{A} B C \bar{D}= \\
& =A \bar{B} \bar{C} D+A B C \bar{D}+A \bar{B} C \bar{D}=A B \bar{C} D+\bar{A} B \bar{C} D+\bar{A} B C \bar{D}
\end{aligned}
$$

(d) $A \bar{B}+A \bar{B} C D+C D+B \bar{C} D+A B C D$

$$
=A \bar{B}(C+\bar{C})(D+\bar{D})+A \bar{B} \bar{C} D+(A+\bar{A})(B+\bar{B}) C D+(A+\bar{A}) B \bar{C} D+A B C D
$$

$$
=A \bar{B} \bar{C} \bar{D}+A \bar{B} \bar{C} D+A B C \bar{D}+A B C D+A \bar{B} \bar{C} D+A B C D+A \bar{B} C D+\bar{A} B C D
$$

$$
+\overline{A B} \underline{C} \underline{D}+A B \bar{C} D+\bar{A} B \underline{C} \bar{C} D+\underset{A}{A B C D}
$$

$=A \bar{B} \bar{C} \bar{D}+A \bar{B} \bar{C} D+A B C \bar{D}+A B C D+A \bar{B} C D+\bar{A} B C D+\overline{A B} C D+A B \bar{C} D+\bar{A} B \bar{C} D$
$=\bar{A} \bar{B} C D+\bar{A} B \bar{C} D+\bar{A} B C D+A \bar{B} C \bar{D}+A \bar{B} \bar{C} D+A \bar{B} C D+A B \bar{C} D+A B C D+A B C \bar{D}$

Reduce the function specified in truth Table to its minimum SOP form by using a Karnaugh map.

S(1). Plot the 1's from table in the text on the map as shown in Fig. and simplify.


$$
X=\bar{B}+C
$$

Use the Karnaugh map method to implement the minimum SOP expression for the logic function specified in truth Table.

| Inputs |  |  |  | Output |
| :---: | :---: | :---: | :---: | :---: |
| $\boldsymbol{A}$ | $\boldsymbol{B}$ | $\boldsymbol{C}$ | $\boldsymbol{D}$ | $\boldsymbol{X}$ |
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 |



## The end of Lecture_07. chapter 4

